

Multicasting in Adhoc Networks Using MAODV Protocol

ALI MUSTAFA

SAAD ZAHID

COMSATS Institute of Information Technology

Attock, Pakistan

AHMED QADDUS

AAMIR ALI RAZA

Bahria University, Islamabad

Pakistan

Abstract:

This paper represent a MAODV (Multicast Ad-hoc On-Demand Distance Vector Protocol) as a multicast routing protocol, it is an extension of AODV which is used for unicasting only. In this research we have evaluated its performance and multicast tree information. Our results show the change in delivery ratio, delay and the overhead in the implementation of this protocol.

While implementing MAODV protocol we have used the environment of FEDORA 8 and a network simulator version NS 2.26 and at the same time it is also implemented and tested in the environment of window vista using CYGWIN.

Key words: ADHOC NETWORKS, MANETS, AODV, MAODV, MULTICASTING

1. Introduction

Mobile adhoc network is the network of mobile nodes which are self organized into arbitrary and having different network topologies allowing peoples to intercommunicate in the area

without and pre-existing infrastructure [1].

In wider area adhoc networks are multi-hop wireless networks and there are many issues which are still challenge and need to be solved (e.g. location management, mobility, security, integrity and etc)[1][2]. There are different types of adhoc networks available according to the different ranges like BAN, PAN, LAN and WAN [1][3].

Adhoc networks are very much appealing to a big extent because of the challenge of maintaining a communication between source and destination [2]. Even in the hard condition when the intermediary node are untrustworthy in participating data in the forward route and swap by the nodes all along the other path [1,3]. So it is not easy to a keep the single source and multiple destinations for data transmission. The known significance of multicast as a mean to decrease the bandwidth utilization for mass distribution of data, and the serious need to communicate in short supply bandwidth over wireless media, it is normal that multicast routing should receive some attention for ad hoc networks [8].

Multicasting in the environments of Manets is the most challenging task because networks is the collection of dynamic nodes which are rapidly changing their position and the same time keep changing their topologies[2]. We can say topologies are temporary in it and the nodes are connected with each other through a very unreliable low bandwidth links[3].

The meaning of reliability varies application to application but when we are talking about mobile adhoc networks then we means all data packets are delivered[2,3]. the sending order of the data packets are maintained and the goal of achieving total number of data packet is achieved[3].

In my simulations and results I will try my best to that all packets should be delivered but it is very challenging in Manets and guaranteeing absolute reliability is not a realistic goal because in mobile networks the receivers and senders are always in dynamic mode and some time they

disconnected from the network for unpredictable slice of time[4]. The only thing on which we can concentrate in Manets to achieve reliability is the maximum packet delivery ratio [3,4].

As I explained above that reliable multicasting is a big challenge in mobile adhoc networks because of the dynamic characteristic of the network only structuring a multicast distribution wouldn't be enough for achieving the desired results and output[5].we need to adopt a alternative path with it which includes the Broadcasting and Dedicated unicast communication for the network[5,6].

The multicast protocols make paths to other hosts on demand. This idea is based on a query-response mechanism same to reactive unicast routing protocols. In the query phase, a node explores the whole environment [4]. Once the query reaches to its destination, the response phase is entered and establishes the path for the multicasting [3].

There comes two structures of multicasting ,the tree based multicasting and mesh based multicasting ,the previous studies and research showed that mesh implementation is much better than the tree structure in generating the results[3][4].

in performance evolution but on the other hands there are many other factors which make mesh structure much complex to implement like to manage the mesh structure of nodes, nodes mobility, time out for unpredictable time, structure less and etc[6].In my research I find that maximum PDR is a challenging task in Manets but my research and simulation improve the performance of tree based structure by using the broadcast of data and by varying the number of senders and the receivers [5].it is very difficult to evaluate the exact number of packets received by receiver and how many packets are lost[5,6].

2. Simulation Environment

Parameters Set in Implementation of MAODV:

- Number of Nodes: Nodes varies from 20 to 50.
- Simulation Area: Area Varies from 300*150 to 1500*300
- Number of Scenarios: 7
- Simulation Time: 300 to 900 Seconds
- Node movement Speed: 0 m/s, 1 m/s and 20 m/s.
- Packet Size: 256
- Traffic: 1, 5, 10, 20 and 50 Kbytes/sec.
- Antenna Range: 100, 150, 200, 250, 500 m

2.1. Scenarios and CBR:

There are total 7 scenarios which we tested in the implementation of MAODV and set the following parameters in it

- Nodes = 20 to 50
- Pause Time = 0.00
- Max speed = 1.00 to 2.00
- Max x = 300.00 to 1500.00
- Max y = 150.00 to 300.00

And about the CBR I have made twenty different combinations by varying number of sender and receiver.

For example

- 1 sender ----- 20 receivers
- 2 senders ----- 40 receivers
- 5 senders ----- 10 receivers
-
- 10 senders ----- 30 receivers

And all these combinations of CBR are used in the scenarios file and all these scenarios files are call in the main TCL file one by one using the different combinations of CBRs.

My simulation atmosphere consists of 50 mobile nodes wandering in a 1500 meters x 300 meters space for 300 seconds of imitation time. The radio communication range is 250 meters. Scenario files are used to determine which nodes are act as receivers or sources and decide when they join or leave a group. In the beginning of the simulation a multicast member node joins the multicast group (approx first 30 seconds) and remains as a member till the whole simulation. Multicast sources start and stop sending packets in the same fashion but using dissimilar arbitrarily generated mobility scenarios. I use only one multicast group for the entire experiments.

Every mobile node moves arbitrarily at a fixed average speed according to a “random waypoint model”. Each node begins its movement from an arbitrary location to a arbitrary destination with a randomly selected speed. When the target is reached, a different random destination is targeted after a break. If we change the pause time it way effect the relative speeds of the mobile nodes. In our experiments we set the pause time to zero to make a harsher mobility situation.

To investigate the protocol MAODV, we perform couples of experiments to discover the performance and behavior of MAODV with changing in different parameters: node mobility, number of senders and receivers, multicast group size and etc etc.

2.2. Number of Senders

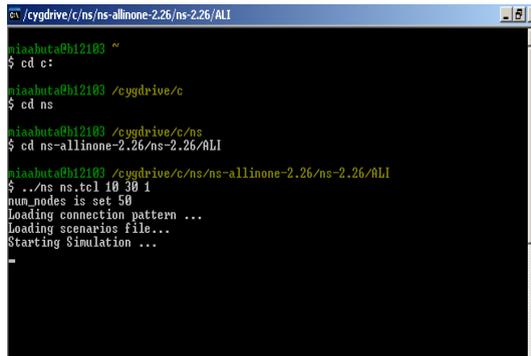
I changed the number of sender in the multicast group to examine the protocol scalability to source nodes and the effect results in traffic load. MAODV, maintain just one group leader for the multicast group that send periodic Group Hellos in the network. In this way, it is extra scalable.

MAODV protocol relies only on a single path on its multicast tree, and must respond to out of order links, by initiate maintenance and repair.

3. Results & Simulation

Simulation of MAODV:

Figure 1 is a simulation representation of MAODV with scenario number and total number of sender and receiver.



```
cydrive/c/ns/ns-allinone-2.26/ns-2.26/ALI
nsaabh121@3 ~
$ cd c:
nsaabh121@3 /cydrive/c
$ cd ns
nsaabh121@3 /cydrive/c/ns
$ cd ns-allinone-2.26/ns-2.26/ALI
nsaabh121@3 /cydrive/c/ns/ns-allinone-2.26/ns-2.26/ALI
$ ./ns ns.tc1 10 30 1
num_nodes is set 50
Loading connection pattern ...
Loading scenarios file ...
Starting Simulation ...
```

Figure 1: Simulation of Scenario 1 in cygwin

In the above diagram it is clearly visible that total numbers of senders are 10 nodes and total receiving nodes are 30 and at that time scenario 1 is running and generating simulation results in trace file. Similarly every scenario run in the same manner and generates its trace file which is further converted in to graph to analyze the behavior of every node at different instant of time.

After simulating the MAODV multiple scripts and scenarios we get the results in table 1. Which clearly shows the packet delivery ratio & latency. Also it gives us information about the total number of send and receives data packets to the destination.

	Receiver	Packet Sent	Packet Received	PDR	Latency
Scenario1	10	5418	45706	0.84	0.036
Scenario2	10	5426	41840	0.77	0.043
Scenario3	10	5383	46254	0.85	0.035
Scenario4	10	5415	41101	0.75	0.038

Scenario5	10	5410	39084	0.72	0.033
Scenario6	10	5425	44718	0.82	0.040
Scenario7	10	5429	32622	0.60	0.030

Table 1: Final results of MAODV simulation.

In these scenarios, I evaluate the performance of MAODV for different traffic: 1, 5, 10, 20 and 50 Kbytes/sec. I calculate the PDR (Packet Delivery Ratio) and the Latency for it. PDR is the ratio of the number of packets sent to the number of packets received and shows the consistency of the protocol. Latency is the measure of average end-to-end packet delay.

3.1 NAM Presentation of MAODV:

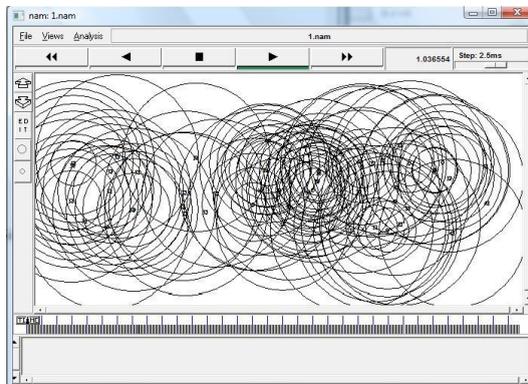


Figure 2: Heavy traffic passing in a MAODV amongst the nodes

As we now that the Manets by nature exhibits dynamic behavior which results different types of region[6,8]. Normally this classification is based upon the density of user. More the user per square in specific area more traffic will be generated in this area. There are two possibilities regarding with the traffic in the region[6,7,8]. The region which is having more traffic is known to be dense region. The traffic which has less traffic is called sparse region [7,8]. These regions come in to existence in a very dynamic fashion. Figure 2, 3, 4 shows the region with different traffic loads within the networks.

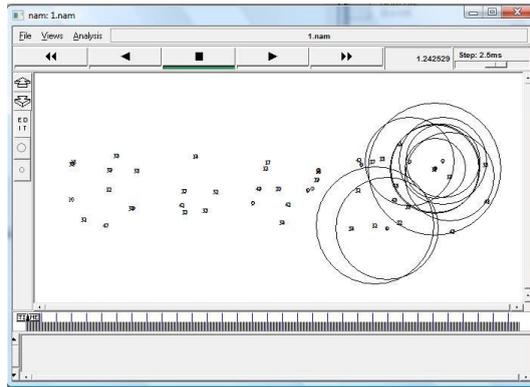


Figure 3: Medium traffic flowing in MAODV multicasting

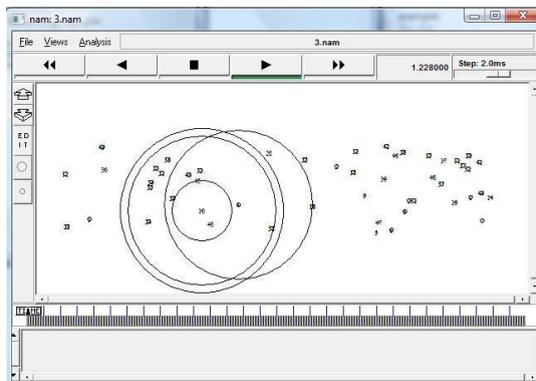


Figure 4: Light Traffic in the multicast network

3.2 Graphical Representation of Different Parameters in MAODV

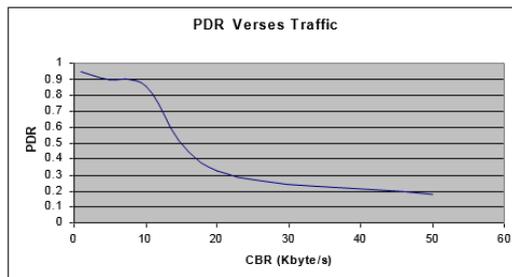


Figure 5: PDR vs. CBR

Fig. 5 shows the rigorous PDR downfall as the traffic increases in MAODV.

When the packets keep on generating at some stage the packet delivery ratio become declining. As the time goes the amount of the data packets is keep on increasing which creates the bottleneck in network which is the big problem in the Manets environment.

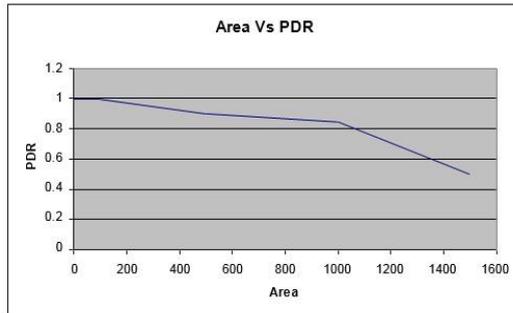


Figure 6: Area vs. PDR

Fig. 6 shows that the PDR of MAODV is healthier for small areas up to 1000.

As the node in the mobile adhoc network is moving from on place to another in highly dynamic manner. The distance among the nodes in increases which take longer time to reach the packets to the destination .therefore most of nodes experience lower packet delivery ratio that is “number of packets that received on destination.

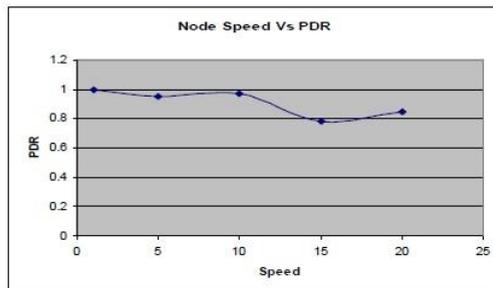


Figure 7: Speed vs. PDR

Fig. 7 shows that the PDR of MAODV is improved for node pace up to 10m/sec.

Speed play very vital role in Manets that is why we call Manets as dynamic networks. as the nodes is moving in very high speed with respect to another node, eventually the distance between the node increase in a drastically manner. therefore most of nodes experience lower packet delivery ratio that is “number of packets that received on destination”.this concept follow the well know relation $s=vt$. Where s is the distance, v is the velocity and t is time.

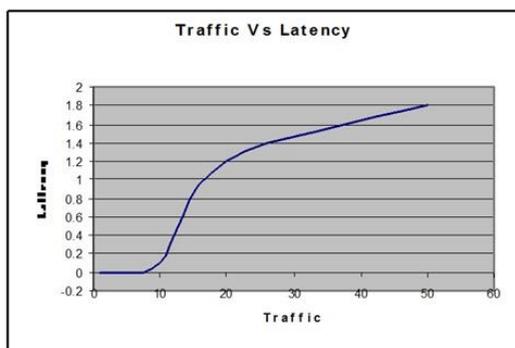


Figure 8: Traffic vs. Latency

Fig. 8 indicates that MAODV protocol has tremendously small latency for traffic up to 10 Kbytes/sec and then for heavier traffic MAODV outperforms.

As the traffic increases the nodes experience the delay because wireless is shared medium. It supports well up to some extent because every device (node) get separate channel. after every channel allocate to every devices that if the new nodes want to send information to another node it will wait until any of the channel will release therefore the delay in communication is produced .the above graph depicts the same scenario which is implemented in the same environment motion above.

In the end of simulation we come to the conclusion that MAODV perform appreciably better. The MAODV performance

is mostly far above the ground under low mobility for only some (1 or 2 senders), consequential in PDRs of approximately 98% and latencies of a hardly any tens of milliseconds. By exploring the trace files in further detail, I come to know that the sole shared tree, with all the operation to keep its hard state, become the log jam and a bottle neck in our all high mobility scenario and for a big amount of multicast sender.

By observing at the network effectiveness, scenarios with comparatively few multicast receivers produce more overhead than scenarios with various multicast receivers, so the charge of structuring and maintaining the multicast tree is amortized over an amplified quantity of packets delivery. Though by using only one shared multicast tree being built and maintained, the overhead grows comparatively slow as we increase the amount of multicast senders.

4. Conclusion

Multicasting can well supported by an extensive range of application that is categorized by a narrow scale of association, classifiably for several MANET applications now envision. Wired network already consists of; entrenched routing protocols exist to offer capable multicasting package. As in different networks nodes become progressively more mobile, these protocols need to develop to give similarly competent service in the new-fangled environment. In Adaptation of wired multicast protocols to MANETs, which are totally missing in infrastructure, appear less promising for the reliability. These existing protocols are designed and planned for permanent networks, may be unsuccessful to keep up with nodes movement and numerous topologies change because of host mobility raise the protocol overheads significantly. To a certain extent, fresh protocols that work in an on-demand approach are being planned and investigate. My investigation

result shows that tree-based on-demand protocol like MAODV lacks in availability of multiple paths

After examining and concluding the experimental results of implementation of different scenarios of MAODV I find that it performs better than any other protocol for the traffic over 10KB/sec.

While talking about the simulation area size MAODV achieved better PDR for the smaller area but if the area size increases it affects the PDR ratio and some time it drops the ratio but in mean while it achieves the best time delay and latency.

4.1. Future Work

We can improve the performance of MAODV by different ways like its tree based on a hard structure, require explicit control messages for the maintenance purposes, all data is flowing through a shared tree and have only a single path for unicasting by the destination, so if we overcome the above mentioned two issues, hard structure and a shared tree it will amongst one of the best ever protocols for multicasting because of its simplicity, easy implementation and reliability of data sending.

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